

COUNTING LYNDON FACTORS

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ABSTRACT. In this paper, we determine the maximum number of distinct Lyndon factors that a word of length n can contain. We also derive formulas for the expected total number of Lyndon factors in a word of length n on an alphabet of size σ , as well as the expected number of distinct Lyndon factors in such a word. The minimum number of distinct Lyndon factors in a word of length n is 1 and the minimum total number is n , with both bounds being achieved by x^n where x is a letter. A more interesting question to ask is *what is the minimum number of distinct Lyndon factors in a Lyndon word of length n ?* In this direction, it is known [13] that an optimal lower bound for the number of distinct Lyndon factors in a Lyndon word of length n is $\lceil \log_\phi(n) + 1 \rceil$, where ϕ denotes the *golden ratio* $(1 + \sqrt{5})/2$. Moreover, this lower bound is attained by the so-called *finite Fibonacci Lyndon words*, which are precisely the Lyndon factors of the well-known *infinite Fibonacci word* f — a special example of a *infinite Sturmian word*. Saari [13] conjectured that if w is Lyndon word of length n , $n \neq 6$, containing the least number of distinct Lyndon factors over all Lyndon words of the same length, then w is a *Christoffel word* (i.e., a Lyndon factor of an infinite Sturmian word). We give a counterexample to this conjecture. Furthermore, we generalise Saari's result on the number of distinct Lyndon factors of a Fibonacci Lyndon word by determining the number of distinct Lyndon factors of a given Christoffel word. We end with two open problems.

1. INTRODUCTION

This paper is concerned with counting Lyndon words occurring in a given word of length n .

First, let us recall some terminology and notation from combinatorics on words (see, e.g., [10, 11]). A *word* is a (possibly empty) finite or infinite sequence of symbols, called *letters*, drawn from a given finite set Σ , called an *alphabet*, of size $\sigma = |\Sigma|$. A finite word $w := x_1x_2 \cdots x_n$ with each $x_i \in \Sigma$ is said to have *length* n , written $|w| = n$. The *empty word* is the unique word of length 0, denoted by ε . The set of all finite words over Σ (including the empty word) is denoted by Σ^* , and for each integer $n \geq 2$, the set of all words of length n over Σ is denoted by Σ^n .

A finite word z is said to be a *factor* of a given finite word w if there exist words u, v such that $w = uzv$. If $u = \varepsilon$, then z is said to be a *prefix* of w , and if $v = \varepsilon$, then z is said to be a *suffix* of w . If both u and v are non-empty, we say that z is a *proper factor* of w . A prefix (respectively, suffix) of w that is not equal to w itself is said to be a *proper prefix* (respectively, *proper suffix*) of w . A factor of an infinite word is a finite word that occurs within it.

A non-empty word x that is both a proper prefix and a proper suffix of a finite word w is said to be a *border* of w . We say that a word which has only an empty border is *borderless*. If, for some word x , $w = xx \cdots x$ (k times for some integer $k \geq 1$), we write $w = x^k$, and w is called the *k -th power* of x . A non-empty finite word is said to be *primitive* if it is not a power of a shorter word. Two finite words u, v are said to be *conjugate* if there exist words x, y such that $u = xy$ and $v = yx$. Accordingly, conjugate words are cyclic shifts of one another, and thus conjugacy is an equivalence relation. A primitive word of length n has exactly n distinct conjugates. For example, the primitive word *abacaba* of length 7 has 7 distinct conjugates; namely, itself and the six words *bacabaa*, *acabaab*, *cabaaba*, *abaabac*, *baabaca*, *aabacab*. The set of all conjugates of a finite word w is called the *conjugacy class* of w .

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In this paper we consider only words on an ordered alphabet $\Sigma = \{a_1, a_2, \dots, a_\sigma\}$ where $a_1 < a_2 < \dots < a_\sigma$. This total order on Σ naturally induces a *lexicographical order* (i.e., an alphabetical order) on the set of all finite words over Σ . A *Lyndon word* over Σ is a non-empty primitive word that is the lexicographically least word in its conjugacy class, i.e., $w \in \Sigma$ or $w < vu$ for all non-empty words u, v such that $w = uv$ (e.g., see [10]). Equivalently, a non-empty finite word w over Σ is Lyndon if and only if $w \in \Sigma$ or $w < v$ for all proper suffixes v of w [7]. Note, in particular, that there is a unique Lyndon word in the conjugacy class of any given primitive word. For example, *aabacab* is the unique Lyndon conjugate of the primitive word *abacaba*. Lyndon words are named after R.C. Lyndon [12], who introduced them in 1954 under the name of “standard lexicographic sequences”. Such words are well known to be borderless [7].

We begin in Section 2 by computing $D(\sigma, n)$, the maximum number of distinct Lyndon factors in a word of length n on an alphabet Σ of size σ . In Section 3 we compute $ET(\sigma, n)$, the expected total number of Lyndon factors (that is, counted according to their multiplicity) in a word of length n over Σ , while Section 4 computes $ED(\sigma, n)$, the expected number of distinct Lyndon factors in word of length n over Σ . Section 5 considers distinct Lyndon factors in a Lyndon word of length n ; in particular, we generalise a result of Saari [13] on the number of distinct Lyndon factors of a *Fibonacci Lyndon word* by determining the number of distinct Lyndon factors of a given *Christoffel word* (i.e., a Lyndon factor of an infinite Sturmian word — to be defined later). Lastly, in Section 6, we state some open problems.

2. THE MAXIMUM NUMBER OF DISTINCT LYNDON FACTORS IN A WORD

Let $D(\sigma, n)$ be the maximum number of distinct Lyndon factors in a word of length n on the alphabet $\Sigma = \{a_1, a_2, \dots, a_\sigma\}$. We want to find a word that achieves $D(\sigma, n)$, given σ and n . It is clear that a necessary condition for attaining the maximum is that w takes the form $a_1^{k_1} a_2^{k_2} \dots a_\sigma^{k_\sigma}$. This word contains $\binom{n+1}{2}$ factors of lengths $1, 2, \dots, n$, of which each is a Lyndon word except those of the form a_i^k , $k > 1$. The number of powers of each a_i is $\binom{k_i+1}{2}$, including a_i itself. The total number of Lyndon factors in w is therefore

$$(1) \quad \binom{n+1}{2} - \sum_{i=1}^{\sigma} \binom{k_i+1}{2} + \sigma,$$

where the final σ counts the single letters a_i . We claim that the summation is minimised when the k_i differ by at most one. Suppose to the contrary that $k_j = k_i + s$ for some i, j and $s \geq 2$. It is easily checked that

$$\binom{k_i+1+s}{2} + \binom{k_i+1}{2} > \binom{k_i+s}{2} + \binom{k_i+2}{2}$$

for $s \geq 2$. Thus the summation term will be minimised when each k_i equals either $\lfloor n/\sigma \rfloor$ or $\lceil n/\sigma \rceil$. If $n = m\sigma + p$, where $0 < p < \sigma$, then $\lfloor n/\sigma \rfloor = m$ and $\lceil n/\sigma \rceil = m+1$. If $p = 0$ then each k_i equals m . We therefore have the following result.

Theorem 1. *If $n = m\sigma + p$, where $0 \leq p < \sigma$, then*

$$(2) \quad D(\sigma, n) = \binom{n+1}{2} - (\sigma - p) \binom{m+1}{2} - p \binom{m+2}{2} + \sigma$$

and the maximum is attained using

$$w = a_1^m \dots a_{n-p}^m a_{n-p+1}^{m+1} \dots a_n^{m+1}.$$

Corollary 2. *If $n = m\sigma$ then*

$$D(\sigma, n) = \binom{\sigma}{2} m^2 + \sigma.$$

Proof. If $n = m\sigma$, Theorem 1 gives

$$\begin{aligned} D(\sigma, n) &= \binom{n+1}{2} - \sigma \binom{m+1}{2} + \sigma \\ &= \frac{\sigma}{2}(m(m\sigma+1) - (m+1)m) + \sigma \\ &= \binom{\sigma}{2}m^2 + \sigma \quad \text{as required.} \end{aligned}$$

□

The following table shows values of $D(\sigma, n)$ for low values of n .

| n | $D(2, n)$ | $D(5, n)$ | $D(10, n)$ |
|-----|-----------|-----------|------------|
| 1 | 2 | 5 | 10 |
| 2 | 3 | 6 | 11 |
| 3 | 4 | 8 | 13 |
| 4 | 6 | 11 | 16 |
| 5 | 8 | 15 | 20 |
| 6 | 11 | 19 | 25 |
| 7 | 14 | 24 | 31 |
| 8 | 18 | 30 | 38 |
| 9 | 22 | 37 | 46 |
| 10 | 27 | 45 | 55 |
| 15 | 58 | 95 | 110 |
| 20 | 102 | 165 | 190 |
| 25 | 158 | 255 | 290 |
| 30 | 227 | 365 | 415 |

TABLE 1. The maximum number of distinct Lyndon factors that can appear in words of length n .

3. THE EXPECTED TOTAL NUMBER OF LYNDON FACTORS IN A WORD

We now wish to calculate the total number $M(\sigma, n)$ of Lyndon factors (that is, counted according to multiplicity) appearing in all words in Σ^n . Consider a Lyndon word L of length $m \leq n$ and a position i , $1 \leq i \leq n - m + 1$, in words of length n . Words containing L starting at position i have the form xLy where xy is any word on Σ with length $n - m$. Thus there will be σ^{n-m} words in Σ^n which contain L in this position. This will be the same for any of the $n - m + 1$ possible values of i so in the σ^n words in Σ^n there will be $(n - m + 1)\sigma^{n-m}$ appearances of L . This is the same for all Lyndon words of this length. The number of such Lyndon words is $1/m$ of the number of primitive words of this length, since exactly one conjugate of each primitive word is Lyndon. The number of primitive words of length n ([10], equation (1.3.7)) is

$$\sum_{d|m} \mu\left(\frac{m}{d}\right) \sigma^d$$

where μ is the Möbius function. To get the total number of Lyndon factors appearing in Σ^n , we sum over possible values of m :

$$(3) \quad M(\sigma, n) = \sum_{m=1}^n \frac{n-m+1}{m} \sigma^{n-m} \sum_{d|m} \mu\left(\frac{m}{d}\right) \sigma^d.$$

Dividing by σ^n gives the expected total number $ET(\sigma, n) := M(\sigma, n)/\sigma^n$ of Lyndon factors in a word of length n on the alphabet Σ . Table 2 below shows values for $\sigma = 2, 5$ and low values of n .

| n | $M(2, n)$ | $ET(2, n)$ | $M(5, n)$ | $ET(5, n)$ |
|-----|-----------|------------|-----------|------------|
| 1 | 2 | 1.00 | 5 | 1.00 |
| 2 | 9 | 2.25 | 60 | 2.40 |
| 3 | 30 | 3.75 | 515 | 4.12 |
| 4 | 87 | 5.43 | 3800 | 6.08 |
| 5 | 234 | 7.31 | 25749 | 8.24 |
| 6 | 597 | 9.32 | 165070 | 10.56 |
| 7 | 1470 | 11.48 | 1018135 | 13.03 |
| 8 | 3522 | 13.76 | 6103350 | 15.62 |
| 9 | 8264 | 16.14 | 35797125 | 18.33 |
| 10 | 19067 | 18.62 | 206363748 | 21.13 |

TABLE 2. Values of the total number $M(\sigma, n)$ of Lyndon factors appearing in all words of length n and the expected total number $ET(\sigma, n)$ of Lyndon factors in a word of length n on an alphabet of size σ for $\sigma = 2, 5$ and $n = 1, 2, \dots, 10$.

4. THE EXPECTED NUMBER OF DISTINCT LYNDON FACTORS IN A WORD

We use the notation from above, with $[n]$ being the set $\{1, 2, \dots, n\}$. Most of the following analysis counts the number of words in Σ^n that contain at least one factor equal to a specific Lyndon word L . At the end we sum over all possible L . Let S be a non-empty set of positions in a word w and let $P(L, S, w) = 1$ if w contains factors equal to L at each position in w beginning at a position in the set S , and 0 otherwise. Note that w may contain other factors equal to L . We claim that

$$(4) \quad \sum_{s=1}^n (-1)^{s+1} \sum_{S \subseteq [n], |S|=s} P(L, S, w) = \begin{cases} 1 & \text{if } w \text{ contains at least one factor equal to } L, \\ 0 & \text{otherwise.} \end{cases}$$

If w contains no factor equal to L then $P(L, S, w)$ equals 0 for all S so the “otherwise” part of the claim holds. Suppose w contains copies of L beginning at positions in $T = \{i_1, i_2, \dots, i_t\}$ and nowhere else. Then $P(L, S, w)$ equals 1 if and only if S is any non-empty subset of T , so the left hand side of (4) becomes

$$\begin{aligned} & \sum_{s=1}^t (-1)^{s+1} |\{S \subseteq T : |S| = s\}| \\ &= \sum_{s=1}^t (-1)^{s+1} \binom{t}{s} \\ &= \sum_{s=0}^t (-1)^{s+1} \binom{t}{s} + 1. \end{aligned}$$

This equals 1 since the final sum is the binomial expansion of $(1 - 1)^t$. The number of words in Σ^n which contain at least one factor equal to L is therefore

$$\begin{aligned} & \sum_{w \in \Sigma^n} \sum_{s=1}^n (-1)^{s+1} \sum_{S \subseteq [n], |S|=s} P(L, S, w) \\ (5) \quad &= \sum_{s=1}^n (-1)^{s+1} \sum_{S \subseteq [n], |S|=s} \sum_{w \in A^n} P(L, S, w). \end{aligned}$$

We now evaluate $\sum_{w \in \Sigma^n} P(L, S, w)$. This is counting the words in Σ^n which have factors L beginning at positions $i \in S$. It clearly equals 0 if $s|L| > n$ since then there is no room in w for s factors L (recalling

that L is Lyndon, therefore borderless, and therefore cannot intersect a copy of itself). We also need the members of S to be separated by at least $|L|$. The number of such sets S is

$$\binom{n - s|L| + s}{s}.$$

Once S is chosen there are $\sigma^{n-s|L|}$ ways of choosing the letters in w which are not in the specified factors L . Thus

$$\sum_{w \in \Sigma^n} P(L, S, w) = \binom{n - s|L| + s}{s} \sigma^{n-s|L|}.$$

Substituting in (5) we see that the number of words in Σ^n which contain at least one occurrence of L is

$$\sum_{s=1}^{\lfloor n/|L| \rfloor} (-1)^{s+1} \binom{n - s|L| + s}{s} \sigma^{n-s|L|}.$$

To get the expected number $ED(\sigma, n)$ of distinct Lyndon factors in a word of length n , we sum this over all L with length at most n , using the same technique as in the previous section, and divide by σ^n . Replacing $|L|$ with m we get the following:

$$(6) \quad ED(\sigma, n) = \sum_{m=1}^n \frac{1}{m} \sum_{d|m} \mu\left(\frac{m}{d}\right) \sigma^d \sum_{s=1}^{\lfloor n/m \rfloor} (-1)^{s+1} \binom{n - sm + s}{s} \sigma^{-sm}.$$

The following table shows values of $ED(\sigma, n)$ for low values of n and several values of σ .

| n | $\sigma = 2$ | 5 | 10 | 20 |
|-----|--------------|-------|-------|-------|
| 1 | 1.00 | 1.00 | 1.00 | 1.00 |
| 2 | 1.75 | 2.20 | 2.35 | 2.42 |
| 3 | 2.50 | 3.56 | 3.94 | 4.14 |
| 4 | 3.25 | 5.02 | 5.69 | 6.05 |
| 5 | 4.06 | 6.55 | 7.57 | 8.12 |
| 6 | 4.91 | 8.16 | 9.54 | 10.31 |
| 7 | 5.81 | 9.82 | 11.59 | 12.61 |
| 8 | 6.77 | 11.54 | 13.70 | 14.99 |
| 9 | 7.77 | 13.31 | 15.88 | 17.45 |
| 10 | 8.83 | 15.13 | 18.11 | 19.97 |
| 15 | 14.77 | 24.93 | 29.90 | 33.36 |
| 20 | 21.67 | 35.76 | 42.58 | 47.70 |
| 25 | 29.35 | 47.43 | 56.02 | 62.73 |
| 30 | 37.70 | 59.82 | 70.11 | 78.33 |

TABLE 3. The expected number $ED(\sigma, n)$ of distinct Lyndon factors in a word of length n for alphabets of size $\sigma = 2, 5, 10, 20$.

5. DISTINCT LYNDON FACTORS IN A LYNDON WORD

Minimising the number of Lyndon factors over words of length n is not very interesting: the minimum number of distinct Lyndon factors is 1 and the minimum total number is n . Both bounds are achieved by x^n where x is a letter. A more interesting question has been studied by Saari [13]: *what is the minimum number of distinct Lyndon factors in a Lyndon word of length n ?* He proved that an optimal lower bound for the number of distinct Lyndon factors in a Lyndon word of length n is

$$\lceil \log_\phi(n) + 1 \rceil$$

where ϕ denotes the *golden ratio* $(1 + \sqrt{5})/2$. Moreover, this lower bound is attained by the so-called finite *Fibonacci Lyndon words*, which are precisely the Lyndon factors of the well-known *infinite Fibonacci word* f — a special example of a *characteristic Sturmian word*.

Following the notation and terminology in [11, Ch. 2], an infinite word s over $\{a, b\}$ is *Sturmian* if and only if there exists an irrational $\alpha \in (0, 1)$, and a real number ρ , such that s is one of the following two infinite words:

$$s_{\alpha, \rho}, s'_{\alpha, \rho} : \mathbb{N} \longrightarrow \{a, b\}$$

defined by

$$s_{\alpha, \rho}[n] = \begin{cases} a & \text{if } \lfloor (n+1)\alpha + \rho \rfloor - \lfloor n\alpha + \rho \rfloor = 0, \\ b & \text{otherwise;} \end{cases} \quad (n \geq 0)$$

$$s'_{\alpha, \rho}[n] = \begin{cases} a & \text{if } \lceil (n+1)\alpha + \rho \rceil - \lceil n\alpha + \rho \rceil = 0, \\ b & \text{otherwise.} \end{cases}$$

The irrational α is called the *slope* of s and ρ is the *intercept*. If $\rho = 0$, we have

$$s_{\alpha, 0} = ac_\alpha \quad \text{and} \quad s'_{\alpha, 0} = bc_\alpha$$

where c_α is called the *characteristic Sturmian word* of slope α . Sturmian words of the same slope have the same set of factors [11, Prop. 2.1.18], so when studying the factors of Sturmian words, it suffices to consider only the characteristic ones.

The *infinite Fibonacci word* f is the characteristic Sturmian word of slope $\alpha = (3 - \sqrt{5})/2$. It can be constructed as the limit of an infinite sequence of so-called *finite Fibonacci words* $\{f_n\}_{n \geq 1}$, defined by:

$$f_{-1} = b, \quad f_0 = a, \quad f_n = f_{n-1}f_{n-2} \quad \text{for } n \geq 1.$$

That is, $f_1 = ab$, $f_2 = aba$, $f_3 = abaab$, $f_4 = abaababa$, $f_5 = abaababaabaab$, etc. (where f_n is a prefix of f_{n+1} for each $n \geq 1$), and we have

$$f = \lim_{n \rightarrow \infty} f_n = abaababaabaab \dots$$

Note. The length of the n -th finite Fibonacci word f_n is the n -th *Fibonacci number* F_n , defined by: $F_{-1} = 1, F_0 = 1, F_n = F_{n-1} + F_{n-2}$ for $n \geq 1$.

More generally, any characteristic Sturmian word can be constructed as the limit of an infinite sequence of finite words. To this end, we recall that every irrational $\alpha \in (0, 1)$ has a unique simple continued fraction expansion:

$$\alpha = [0; a_1, a_2, a_3, \dots] = \frac{1}{a_1 + \frac{1}{a_2 + \frac{1}{a_3 + \dots}}}$$

where each a_i is a positive integer. The n -th *convergent* of α is defined by

$$\frac{p_n}{q_n} = [0; a_1, a_2, \dots, a_n] \quad \text{for all } n \geq 1,$$

where the sequences $\{p_n\}_{n \geq 0}$ and $\{q_n\}_{n \geq 0}$ are given by

$$\begin{aligned} p_0 &= 0, & p_1 &= 1, & p_n &= a_n p_{n-1} + p_{n-2}, & n &\geq 2 \\ q_0 &= 1, & q_1 &= a_1, & q_n &= a_n q_{n-1} + q_{n-2}, & n &\geq 2 \end{aligned}$$

Suppose $\alpha = [0; 1 + d_1, d_2, d_3, \dots]$, with $d_1 \geq 0$ and all other $d_n > 0$. To the *directive sequence* (d_1, d_2, d_3, \dots) , we associate a sequence $\{s_n\}_{n \geq -1}$ of words defined by

$$s_{-1} = b, \quad s_0 = a, \quad s_n = s_{n-1}^{d_n} s_{n-2} \quad \text{for } n \geq 1.$$

Such a sequence of words is called a *standard sequence*, and we have

$$|s_n| = q_n \quad \text{for all } n \geq 0.$$

Note that ab is a suffix of s_{2n-1} and ba is a suffix of s_{2n} for all $n \geq 1$.

Standard sequences are related to characteristic Sturmian words in the following way. Observe that, for any $n \geq 0$, s_n is a prefix of s_{n+1} , which gives obvious meaning to $\lim_{n \rightarrow \infty} s_n$ as an infinite word. In fact, one can prove [8, 3] that each s_n is a prefix of c_α , and we have

$$c_\alpha = \lim_{n \rightarrow \infty} s_n.$$

The following lemma collects together some properties of the *standard words* s_n . Note that from now on when referring to Lyndon words over the alphabet $\{a, b\}$ we assume the natural order $a < b$.

Lemma 3. *Let c_α be the characteristic Sturmian word of slope $\alpha = [0; 1 + d_1, d_2, d_3, \dots]$ with $c_\alpha = \lim_{n \rightarrow \infty} s_n$ where the words s_n are defined as above.*

- For all $n \geq 1$, s_n is a primitive word [6].
- For all $n \geq 1$, there exist uniquely determined palindromes u_n, v_n, p_n such that

$$s_n = u_n v_n = \begin{cases} p_n ab & \text{if } n \text{ is odd,} \\ p_n ba & \text{if } n \text{ is even,} \end{cases}$$

where $|u_n| = q_{n-1} - 2$ and $|v_n| = q_n - q_{n-1} + 2$. [6]

- For all $n \geq 1$, the reversal of s_n is the $(q_n - 2)$ -nd conjugate of s_n , and hence the conjugacy class of s_n is closed under reversal. [9, Prop. 2.9(4)]
- The Lyndon factors of c_α of length at least 2 are precisely the Lyndon conjugates of the (primitive) standard words s_n for all $n \geq 1$. [2, 5]

The following lemma is a generalisation of [13, Lemma 8].

Lemma 4. *Let c_α be the characteristic Sturmian word of slope $\alpha = [0; 1 + d_1, d_2, d_3, \dots]$ with $c_\alpha = \lim_{n \rightarrow \infty} s_n$ where $s_n = p_n xy$ with $xy \in \{ab, ba\}$. The Lyndon conjugate of s_n is the word ap_nb for all $n \geq 1$. Moreover, every Lyndon factor of c_α that is shorter than ap_nb is either a prefix or a suffix of ap_nb .*

Proof. First we show that, for all $n \geq 1$, the Lyndon conjugate of s_n is the word ap_nb . If $s_n = p_n ba$, then ap_nb is clearly a conjugate of s_n and it is Lyndon [2, 5]. On the other hand, if $s_n = p_n ab$, then bp_na is clearly a conjugate of s_n , and since the conjugacy class of s_n is closed under reversal and p_n is a palindrome (by Lemma 3), it follows that ap_nb is a conjugate of s_n and it is Lyndon [2, 5].

To prove the second claim, it suffices to show that if $k < n$, then the Lyndon conjugate of s_k is a prefix or suffix of ap_nb (since, by Lemma 3, the Lyndon factors of c_α of length at least 2 are precisely the Lyndon conjugates of the (primitive) standard words in c_α). The claim is true for $k = -1$ and $k = 0$ since $s_{-1} = b$ and $s_0 = a$. It is also true for $k = 1$ because $s_1 = a^{d_1}b$ is the Lyndon conjugate of itself, and is a prefix of ap_nb if $d_1 \geq 1$ and a suffix of ap_nb if $d_1 = 0$. Now suppose that $k \geq 2$. Then $k < n$ implies that s_k is a prefix of p_n . Furthermore, since p_n is a palindrome, the reversal of s_k is a suffix of p_n . Therefore if $s_k = p_k ba$, then its Lyndon conjugate ap_kb is a prefix of ap_nb ; otherwise, if $s_k = p_k ab$, then its Lyndon conjugate ap_kb is a suffix of ap_nb . \square

The Lyndon factors of (characteristic) Sturmian words of length at least 2 (i.e., the Lyndon conjugates of standard words) over $\{a, b\}$ are precisely the so-called *Christoffel words* beginning with the letter a (see, e.g., the nice survey [1]). Christoffel words take the form $aPal(v)b$ and $bPal(v)a$ where $v \in \{a, b\}^*$ and Pal is iterated palindromic closure, defined by:

$$Pal(\varepsilon) = \varepsilon \quad \text{and} \quad Pal(wx) = (Pal(w)x)^+ \quad \text{for any finite word } w \text{ and letter } x,$$

Proof. The word w is the Lyndon conjugate of the standard word $s_{n+1} = s_n s_{n-1}$ with $s_{-1} = b$, $s_0 = a$, and $s_i = s_{i-1}^{d_i} s_{i-2}$ for $1 \leq i \leq n$ (see Remark 5). By Lemma 4, the Lyndon conjugates of s_i for $1 \leq i \leq n$ are either prefixes or suffixes of w . Moreover, for each i with $1 \leq i \leq n$, the standard word s_i contains d_i distinct Lyndon factors of lengths $|s_{i-1}^m s_{i-2}|$ for $m = 1, 2, \dots, d_i$. By Lemma 3, these are the only Lyndon factors of the Lyndon word w besides itself and the two letters a and b . Hence $\mathcal{L}(w) = (d_1 + d_2 + \dots + d_n) + 3$. \square

The above result is a generalisation of [13, Lemma 9], which reworded (with the indexing of Fibonacci words and numbers shifted back by 2) states that if w is the Fibonacci Lyndon word of length F_{n-2} for some $n \geq 3$, i.e., the Christoffel word of slope

$$F_{n-4}/F_{n-2} = [0; 2, \underbrace{1, 1, \dots, 1}_{(n-3) \text{ 1s}}]$$

beginning with the letter a , then $\mathcal{L}(w) = (n-3) + 3 = n$.

Examples:

- The Christoffel word of slope $2/5 = [0; 2, 1, 1]$ beginning with the letter a , namely $aPal(ab)b = aabab$, is the Fibonacci Lyndon word of length $F_3 = 5$ that contains the minimum number ($5 = 2 + 3$) of distinct Lyndon factors for its length.
- The Christoffel word of slope $1/6 = [0; 5, 1]$ beginning with the letter a , namely $aPal(a^4)b = aaaaab$, is a Sturmian Lyndon word of length 6 containing the minimum number ($7 = 4 + 3$) of distinct Lyndon factors for its length.
- The Christoffel word of slope $3/8 = [0; 2, 1, 1, 1]$ beginning with the letter a , namely $aPal(aba)b = aabaabab$, is the Fibonacci Lyndon word of length $F_4 = 8$ containing the minimum number ($6 = 3 + 3$) of distinct Lyndon factors for its length.

6. OPEN PROBLEMS

Open Problem 1: One might suspect that any given Christoffel word contains the minimum number of distinct Lyndon factors over all Lyndon words of the same length. However, this is not true. For instance, the Christoffel word of slope $5/11 = [0; 2, 4, 1]$ beginning with the letter a , namely $aPal(ab^4)b = aababababab$, contains 8 ($= 5 + 3$) distinct Lyndon factors, but the minimum number of distinct Lyndon factors of a Lyndon word of length 11 is actually 7. This minimum is attained by the Christoffel word $aPal(a^2ba)b = aaabaaabaab$ of slope $3/11 = [0; 3, 1, 1, 1]$.

Is it true that the minimum number of distinct Lyndon factors over all Lyndon words of the same length is attained by at least one Christoffel word of that length?

Open Problem 2: Tables 2 and 3, showing values for $ET(\sigma, n)$ and $ED(\sigma, n)$, raise the question of whether there may exist asymptotic formulas for these quantities, simpler than the exact values displayed in equations (2) and (6), respectively.

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